# **AN14474**

## i.MX 9 - L3 Cache Partitioning for Predictable Real-Time Performance

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#### **Document information**

Information	Content
Keywords	AN14474, L3 Cache, L3 Cache Partitioning, Real-Time, i.MX 93, i.MX 95, ARM DynamIQ Shared Unit, DSU
Abstract	This application note describes how to partition the L3 cache between the cores, using features of the Arm DynamIQ Shared Unit.

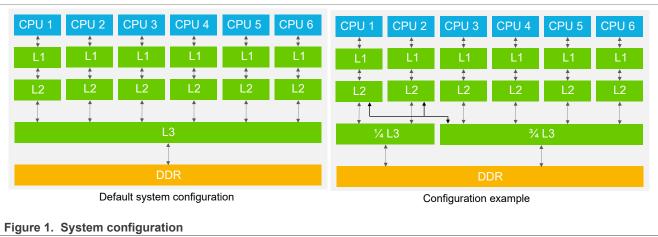


i.MX 9 - L3 Cache Partitioning for Predictable Real-Time Performance

## Introduction

In most CPU clusters, L3 cache is a shared resource, which typically gives the best overall performance of the system for a given cache size. However, this may not be ideal in real-time situations. For example, if a lowpriority task is memory-intensive, it may pollute the entire L3 cache of the cluster, increasing the latency of a higher-priority task. This is undesirable in a real-time environment.

This document describes how to partition the L3 cache between the cores, using features of the Arm DynamIQ Shared Unit. Furthermore, it shows how to allocate specific tasks (high priority real-time) on cores with dedicated partitions of L3 cache. The code in this document was tested on i.MX 95 and i.MX 93.



As shown in Figure 1,

- the left image shows the default L3 cache configuration on i.MX 95, shared among all the cores.
- the right image shows a possible configuration in which \( \frac{1}{4} \) of L3 cache is used exclusively by CPU 1 and CPU 2, while the remaining 3/4 of L3 cache is shared among all six cores.

#### General approach 2

For a complete description of how to partition L3 cache, see Arm® DynamIQ™ Shared Unit Technical Reference Manual. In a nutshell, the L3 cache is divided in four equal way groups (a fancy name for parts), numbered 0-3. Each way group can be assigned to one or more of the eight "schemes", numbered 0-7. A scheme is simply a set of way groups. All unassigned way groups are shared among all eight schemes. Each CPU has to be allocated to one of the schemes and have access to the cache of that scheme. To implement the example in Figure 1, perform the following steps:

- 1. Assign way group 0 of the cache to scheme ID 1.
- 2. Leave way groups 1-3 of the cache unassigned (or, alternatively, assign way groups 1-3 to both schemes ID 0 and ID 1).
- 3. Set CPU 1-2 to use scheme ID 1.
- 4. Set CPU 3-6 to use scheme ID 0.

It is obligatory that all **used** schemes have exclusive access, or shared access to at least one-way group of the cache.

To implement the cache partitioning, this solution provides access from the Linux user-space to the following registers: CLUSTERPARTCR EL1 and CLUSTERTHREADSID EL1. This provides flexibility and ease of configuration, but some protections may be required in a production system. For the complete description of the registers, see the manual referenced above. Briefly:

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- CLUSTERPARTCR EL1 configures the allocation of the cache way groups to the schemes.
- CLUSTERTHREADSID EL1, one per core, allows setting the scheme ID used by that core.

In this implementation, CLUSTERACPSID\_EL1 and CLUSTERSTASHSID\_EL1 are left on their default value (0), which means that ACP transactions and stash requests are directed to scheme ID 0. Make sure that scheme ID 0 has access to at least one way group.

This solution adds a kernel module, which exposes one read/write sysfs file for each of the relevant registers. Writes and reads to these files are forwarded to a patched TF-A, which is able to read/write these registers (running at Exception Level 3).

For testing, we reserve a part of the cache for the exclusive use of some of the cores. We evaluate the performance of a real-time task running on the cores with exclusive cache while the rest of the processes, run on the other cores. We use memory-intensive tasks which are strongly affected by the cache performance.

## 3 Implementation

The current implementation is based on the LF-6.6.23 2.0.0 BSP release. Other versions may require porting.

- On the Linux PC, set up the Yocto environment according to Section 3, 4, and 5, in the i.MX Yocto Project User's Guide (document <u>UG10164</u>). For i.MX93 only you can, alternatively, use the <u>Real-Time Edge</u> <u>Software</u>, according to Section 3, 4 and 5 in the Real-time Edge Yocto Project User Guide (document <u>RTEDGEYOCTOUG</u>).
- 2. Clone the <u>recipes-cachepartition</u> repository in the meta-imx/meta-imx-bsp directory. The recipes-cachepartition directory contains the following recipe appends:
  - A kernel patch implementing the kernel module for the L3 cache partitioning (in linux-imx subdirectory).
  - An ATF patch implementing the necessary SMC calls for setting the registers (in imx-atf).
  - A small tool, usecache, which can be used to test the amount of available cache (in usecache).

```
$ cd ~/imx-yocto-bsp/sources/meta-imx/meta-imx-bsp
$ git clone -b lf-6.6.23-2.0.0 https://github.com/nxp-imx-support/recipes-
cachepartition
```

- 3. To use the Preempt-RT kernel in the standard Linux distribution, see the <u>How to Use the Preempt-RT Kernel</u> in the Standard Yocto Linux BSP.
- 4. rt-tests contains the cyclictest tool, useful to measure the system latency. You can build it by adding the following lines to the conf/local.conf file.

```
CORE_IMAGE_EXTRA_INSTALL += " rt-tests"
```

## Note:

[i.MX 93] If you are using Real-Time Edge, the Preempt-RT kernel is default, and these lines are not needed

5. Add the usecache package to the image. usecache is a cache stress test, which can be used for the L3 cache partitioning validation. Add the below line in the conf/local.conf file.

```
CORE IMAGE EXTRA INSTALL += " usecache"
```

6. Build the imx-image-full image.

```
bitbake imx-image-full
```

#### Note:

[i.MX 93] If you are using Real-Time Edge, run, instead, the following command:

```
bitbake nxp-image-real-time-edge
```

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7. Write the resulted <image\_name>.wic.zst image located in the tmp/deploy/images/<machine> directory on the SD card using the following command:

```
$ zstd -d <image_name>.wic.zst
$ sudo dd if=<image_name>.wic of=/dev/sd<x> bs=1M conv=fsync
```

## 4 Testing

To test, perform the following steps:

- Connect the USB debug port of the board to the PC using a USB cable. This action creates four virtual serial ports on the PC. Typically, the third serial port corresponds to the Linux console. Open this port in a terminal emulator using the following parameters: 115200 baud rate, 8 data bits, no parity, and 1 stop bit.
- 2. Boot the board.
- 3. Load the cachepartition module. The module prints various debug messages in the kernel log, which can be inspected with <code>dmesg</code>.

```
modprobe cachepartition
```

- 4. Go to the /sys/kernel/cachepartition directory. Here locate the following files: PARTCR\_EL1, and THREADSID EL1 [0-5] one for each core.
- 5. You can check the current value of the registers using cat. Example:

```
cat partcr el1
```

6. You can set the value of the registers using echo. Example:

```
echo e1 > partcr el1
```

## 4.1 Examples of configuration

To configure, perform the following steps:

1. Default configuration: All the L3 caches are shared among all cores.

```
echo 0 > partcr_el1
```

2. ¾ L3 cache to cores 4-5 (Scheme ID 1), and ¼ L3 cache to cores 0-3, ACP and STASH (Scheme ID 0). The value written in the partcr\_ell register is E1, which is 11100001 in binary. It means that way groups 3, 2, 1 are assigned to Scheme ID 1 (first 4 bits) and the way group 0 – assigned to Scheme ID 0. Then, we select Scheme ID 0 for cores 0-3 and Scheme ID 1 for cores 4-5.

```
echo e1 > partcr_el1
echo 0 > threadsid_el1_0
echo 0 > threadsid_el1_1
echo 0 > threadsid_el1_2
echo 0 > threadsid_el1_3
echo 1 > threadsid_el1_4
echo 1 > threadsid_el1_5
```

You can check that the configuration is working using the usecache tool.

```
taskset -c 0 usecache 512 100 0 taskset -c 4 usecache 512 100 4
```

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The previous commands run the same usecache test on CPU 0 and on CPU 4. usecache is a memory intensive task, and it uses, in this case, 512 K of memory. Given that CPU 4 has access to more L3 cache than CPU 0, it runs significantly faster.

```
taskset -c 0 usecache 64 100 0 taskset -c 4 usecache 64 100 4
```

In this case, the memory space used by the usecache is 64 K, and fits in the L1-L3 cache of all CPUs, so it runs in approximately equal time.

```
taskset -c 0 usecache 512 100 0 & taskset -c 4 usecache 512 100 4
```

In this case, we run usecache simultaneously on CPU 0 and CPU 4, but because they have access to separate parts of the cache, the time remains about the same as when we run separately.

```
taskset -c 0 usecache 512 100 0 & taskset -c 1 usecache 512 100 1
```

In this case, we run usecache simultaneously on CPU 0 and CPU 1, but because they share the same L3 cache partition, the time increases, compared to when run separately.

3. Cores 4-5 have access to all the L3 cache (Scheme ID 2), cores 0-3 have access only to the first ¼ of L3 cache (Scheme ID 1), ACP and STASH have access only to the second ¼ of L3 cache (Scheme ID 0). The value F84 written in the partcr\_ell register is 111110000100, which means that Scheme ID 2 has access to all 4-way groups (the four most significant bits), Scheme ID 1 has access only to way group 3 (the middle four bits) and Scheme ID 0 has access only to way group 2 (the four least significant bits).

```
echo f84 > partcr_el1
echo 1 > threadsid_el1_0
echo 1 > threadsid_el1_1
echo 1 > threadsid_el1_2
echo 1 > threadsid_el1_3
echo 2 > threadsid_el1_4
echo 2 > threadsid_el1_5
```

Again, you can use usecache to test the performance in this case.

To compile the rt-tests package and the Preempt-RT Linux kernel, use the cyclictest to test the system latency on the various cores.

When using the Jailhouse hypervisor, configure the L3 cache before enabling Jailhouse. Once Jailhouse is enabled, it restricts the communication between the kernel and the ATF. You can always disable Jailhouse temporarily, to change the L3 cache configuration.

## 5 Note about the source code in the document

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## 6 Revision history

Table 1 summarizes the revisions to this document.

Table 1. Revision history

Document ID	Release date	Description
AN14474 v.1.2	19 February 2025	Use the GitHub repository for the recipes-cachepartition, instead of a SW archive
AN14474 v.1.1	28 October 2024	Added hyperlink of AN14474SW in Section 3
AN14474 v.1.0	24 October 2024	Initial public release

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